

Announcements

- MidTerm Exam 1: October 16 in class
- MidTerm Exam 2: Last day of class
- Final: NO FINAL EXAM

QuickSelect: a variant of QuickSort

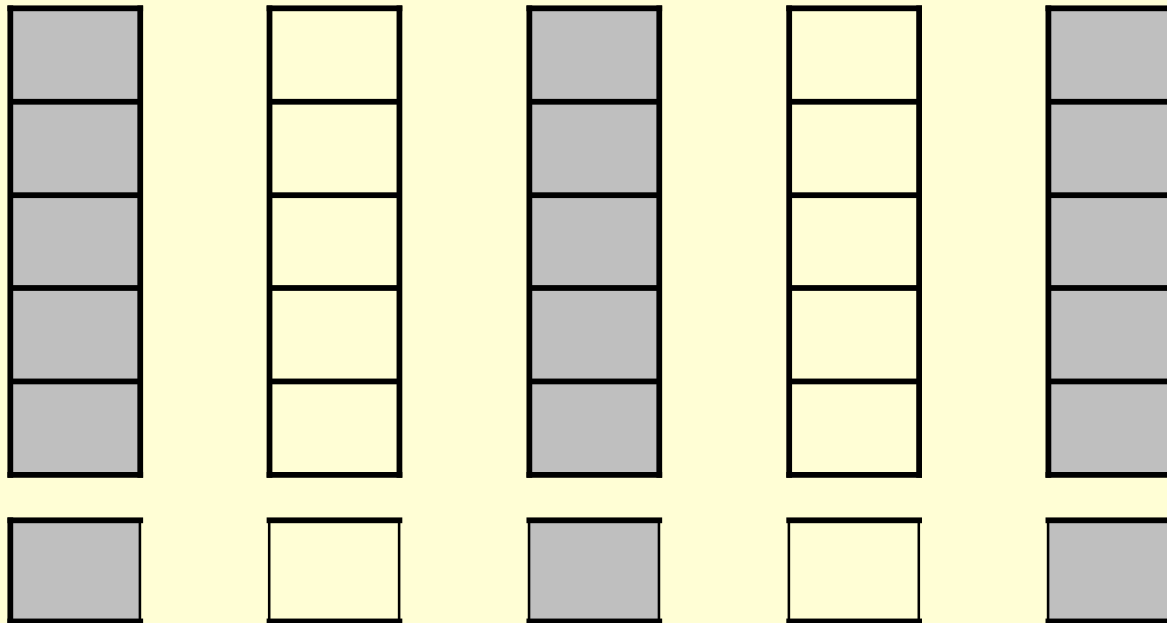
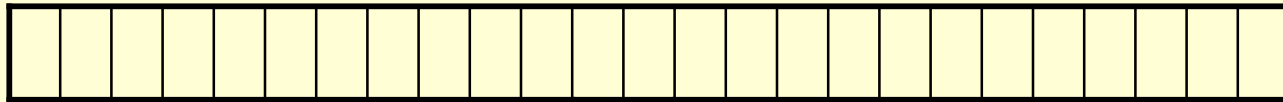
QUICKSELECT(*array A, int k, int p, int r*)

▷ Select k -th largest in subarray $A[p..r]$

```
1  if ( $p = r$ )
2      then return  $A[p]$ 
3   $q \leftarrow$  PARTITION( $A, p, r$ )
4   $i \leftarrow q - p + 1$     ▷ Compute rank of pivot
5  if ( $i = k$ )
6      then return  $A[q]$ 
7  if ( $i > k$ )
8      then return QUICKSELECT( $A, k, p, q - 1$ )
9  else return QUICKSELECT( $A, k - i, q + 1, r$ )
```

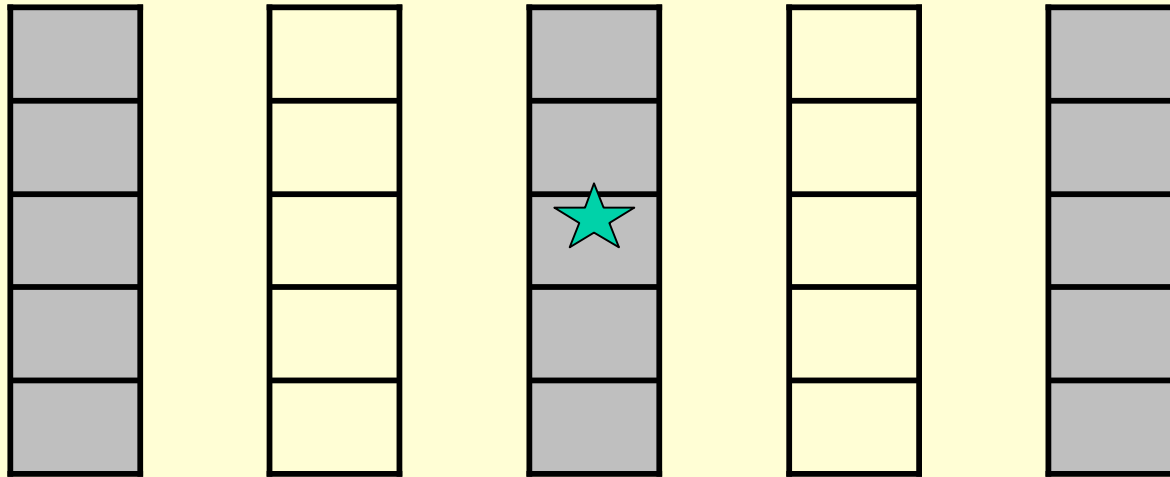
k-Selection & Median: Improved Algorithm

- Start with initial array



k-Selection & Median: Improved Algorithm(Cont'd)

- Use median of medians as pivot



- $T(n) < O(n) + T(n/5) + T(3n/4)$

ImprovedSelect

IMPROVEDSELECT(*array A, int k, int p, int r*)

▷ Select k -th largest in subarray $A[p..r]$

```
1  if ( $p = r$ )
2    then return  $A[p]$ 
3    else  $N \leftarrow r - p + 1$ 
4    Partition  $A[p..r]$  into subsets of 5 elements and
   collect all medians of subsets in  $B[1..\lceil N/5 \rceil]$ .
5     $Pivot \leftarrow$  IMPROVEDSELECT( $B, 1, \lceil N/5 \rceil, \lceil N/10 \rceil$ )
6     $q \leftarrow$  PIVOTPARTITION( $A, p, r, Pivot$ )
7     $i \leftarrow q - p + 1$     ▷ Compute rank of pivot
8    if ( $i = k$ )
9      then return  $A[q]$ 
10   if ( $i > k$ )
11     then return IMPROVEDSELECT( $A, k, p, q - 1$ )
12     else return IMPROVEDSELECT( $A, k - i, q + 1, r$ )
```

PivotPartition

PIVOTPARTITION(*array A*, *int p*, *int r*, *item Pivot*)

▷ Partition using provided *Pivot*

```
1  $i \leftarrow p - 1$ 
2 for  $j \leftarrow p$  to  $r$ 
3     do if ( $A[j] \leq Pivot$ )
4         then  $i \leftarrow i + 1$ 
5             exchange  $A[i] \leftrightarrow A[j]$ 
6 return  $i + 1$ 
```

Analysis of ImprovedSelect

Number of elements greater than “median of medians” is at least

$$3 \left(\left\lceil \frac{1}{2} \left\lceil \frac{n}{5} \right\rceil \right\rceil - 2 \right) \geq \frac{3n}{10} - 6$$

Why?

Our recurrence is given by:

$$T(n) = O(n) + T(\lceil n/5 \rceil) + T(3n/4)$$

Thus there exists a positive constant a such that

$$T(n) \leq an + T(\lceil n/5 \rceil) + T(3n/4)$$

Using the substitution method, let's guess that $T(n) = O(n)$, i.e., $T(n) \leq cn$. Then we need to show that

$$an + c\lceil n/5 \rceil + c(3n/4) \leq cn$$

What positive values of c and n_0 would enforce the above inequality? When $n > 70$, and choosing $c \geq 20a$ will satisfy above inequality.

Data Structure Evolution

- Standard operations on data structures
 - Search
 - Insert
 - Delete
- Linear Lists
 - Implementation: Arrays (Unsorted and Sorted)
- Dynamic Linear Lists
 - Implementation: Linked Lists
- Dynamic Trees
 - Implementation: Binary Search Trees