

Announcements

- MidTerm Exam 1: October 16 in class
- MidTerm Exam 2: Last day of class
- Final: NO FINAL EXAM

Data Structure Evolution

- Standard operations on data structures
 - Search
 - Insert
 - Delete
- Linear Lists
 - Implementation: Arrays (Unsorted and Sorted)
- Dynamic Linear Lists
 - Implementation: Linked Lists
- Dynamic Trees
 - Implementation: Binary Search Trees

BST: Search

TREESearch(*node x*, *key k*)

▷ Search for key k in subtree rooted at node x

1 **if** $((x = \text{NIL}) \text{ or } (k = \text{key}[x]))$

2 **then return** x

3 **if** $(k < \text{key}[x])$

4 **then return** TREESearch($\text{left}[x]$, k)

5 **else return** TREESearch($\text{right}[x]$, k)

Time Complexity: $O(h)$

h = height of binary search tree

Not $O(\log n)$ — Why?

BST: Insert

TREEINSERT(*tree* T , *node* z)

▷ Insert node z in tree T

1 $y \leftarrow \text{NIL}$

2 $x \leftarrow \text{root}[T]$

3 **while** ($x \neq \text{NIL}$)

4 **do** $y \leftarrow x$

5 **if** ($\text{key}[z] < \text{key}[x]$)

6 **then** $x \leftarrow \text{left}[x]$

7 **else** $x \leftarrow \text{right}[x]$

8 $p[z] \leftarrow y$

9 **if** ($y = \text{NIL}$)

10 **then** $\text{root}[T] \leftarrow z$

11 **else if** ($\text{key}[z] < \text{key}[y]$)

12 **then** $\text{left}[y] \leftarrow z$

13 **else** $\text{right}[y] \leftarrow z$

Time Complexity: $O(h)$

h = height of binary search tree

Search for x in T

Insert x as leaf in T

BST: Delete

Time Complexity: $O(h)$

h = height of binary search tree

TREEDeLETE(*tree* T , *node* z)

▷ Delete node z from tree T

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1  if ((left[ $z$ ] = NIL) or (right[ $z$ ] = NIL))
2      then  $y \leftarrow z$ 
3      else  $y \leftarrow \text{TREE-SUCCESSOR}(z)$ 
4  if (left[ $y$ ]  $\neq$  NIL)
5      then  $x \leftarrow \textit{left}[y]$ 
6      else  $x \leftarrow \textit{right}[y]$ 
7  if ( $x \neq \text{NIL}$ )
8      then  $p[x] \leftarrow p[y]$ 
9  if ( $p[y] = \text{NIL}$ )
10     then  $\textit{root}[T] \leftarrow x$ 
11     else if ( $y = \textit{left}[p[y]]$ )
12         then  $\textit{left}[p[y]] \leftarrow x$ 
13         else  $\textit{right}[p[y]] \leftarrow x$ 
14  if ( $y \neq z$ )
15     then  $\textit{key}[z] \leftarrow \textit{key}[y]$ 
16         cop  $y$ 's satellite data into  $z$ 
17  return  $y$ 
```

Set y as the node to be deleted.
It has at most one child, and let
that child be node x

If y has one child, then y is deleted
and the parent pointer of x is fixed.

The child pointers of the parent of x
is fixed.

The contents of node z are fixed.

Animations

- **BST:**

http://babbage.clarku.edu/~achou/cs160/examples/bst_animation/BST-Example.html

- **Rotations:**

http://babbage.clarku.edu/~achou/cs160/examples/bst_animation/index2.html

- **RB-Trees:**

http://babbage.clarku.edu/~achou/cs160/examples/bst_animation/RedBlackTree-Example.html