# Tree Augmentation 

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## The Problem: CodeChef CHN15E

- Given tree $T$, the augmented tree $\mathrm{G}_{\mathrm{T}}$ is defined as the graph obtained by joining every pair of vertices at distance 2 from each other.
- The problem is to construct $T$, given $\mathrm{G}_{\mathrm{T}}$.


## Simple Properties

- Vertices of $T$ and $G_{T}$ are the same.
- Let neighbors of vertex $v$ in $T$ be the set $N(v)$
- The set $\{v\} \cup N(v)$ forms a clique in $G_{T}$.
- A subset of vertices in a graph forms a clique if all of them are connected by edges (i.e., no pair of vertices in this subset are missing an edge)
- A maximal clique is a set of vertices that forms a clique for which no superset is a clique.


## More Properties

- For a tree $T$ with $n$ vertices, the augmented tree $G_{T}$ has at most $n$ maximal cliques
- Each maximal clique of $\mathrm{G}_{\mathrm{T}}$ looks like this:
- \{v\} U N(v)
- There are no other maximal cliques in $\mathrm{G}_{\mathrm{T}}$.
- If tree $T$ is just a star (one vertex connected to all others), then $G_{T}$ is a simple clique
- If $\mathrm{G}_{\mathrm{T}}$ is not a clique, then it has more than one maximal clique, and then $T$ is not a star.


## One more important property

- If $(x, y)$ is an edge of $T$
- Then the vertices $x$ and $y$ appear together in exactly two maximal cliques, except if one of them is a leaf
- If one of them is a leaf, then they appear together in exactly one maximal clique


## Properties of Cliques of $\mathrm{G}_{\mathrm{T}}$

- Vertex $v$ is present in <= deg(v)+1 maximal cliques
- $\operatorname{Deg}(v)$ is degree of vertex $v$
- If $v$ has $k>0$ leaves as neighbors in $T$, then $v$ is present in exactly $\operatorname{deg}(\mathrm{v})-\mathrm{k}+1$ maximal cliques
- If $v$ has $m$ non-leaves as neighbors in $T$, then $v$ is in
- Exactly $m+1$ maximal cliques, if $v$ is not a leaf
- If $v$ has no leaves as neighbors in $T$, then $v$ is in
- exactly $\operatorname{deg}(\mathrm{v})+1$ maximal cliques, if v is not a leaf
- If $v$ is a leaf, it is in exactly 1 maximal clique


## Algorithmic Ideas

1. Identify all maximal cliques of $G_{T}$
2. For each vertex v, compute

- C $[v]=$ \# of maximal cliques of $G_{T}$ containing $v$

3. Identify leaves of $T$ : all vertices with $C[v]=1$
4. Figure out how many non-leaf neighbors each vertex has.
5. Figure out pairs of non-leaf vertices connected by an edge (present in exactly 2 max cliques)

## More Properties of leaves of T

- If two leaves $x$ and $y$ are connected to the same non-leaf node, then they appear together in exactly one maximal clique and in no other clique
- If two leaves $x$ and $y$ are not connected to the same non-leaf node, then they never appear together in a maximal clique


## Algorithmic Ideas

1. Figure out all leaves of $T$
2. Identify all edges of T connecting non-leaves (skeleton $\mathrm{T}^{\prime}$ )
3. Figure out groups of leaves connected to same non-leaf
4. Figure out which leaf is connected to which non-leaf:
a) Construct skeleton $T^{\prime}$
b) Construct maximal cliques of $\mathrm{T}^{\prime}$ corresponding to non-leaf
c) Each maximal clique $A^{\prime}$ of $T^{\prime}$ corresponds to only one maximal clique $A$ of $G_{T}$ and to one non-leaf node $v$.
d) Connect all leaf nodes in A to non-leaf node $v$
