#### **Tree Augmentation**

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## The Problem: CodeChef CHN15E

- Given tree T, the augmented tree G<sub>T</sub> is defined as the graph obtained by joining every pair of vertices at distance 2 from each other.
- The problem is to construct T, given  $G_T$ .

## **Simple Properties**

- Vertices of T and  $G_T$  are the same.
- Let neighbors of vertex v in T be the set N(v)
- The set {v} U N(v) forms a clique in G<sub>T</sub>.
  - A subset of vertices in a graph forms a clique if all of them are connected by edges (i.e., no pair of vertices in this subset are missing an edge)
- A maximal clique is a set of vertices that forms a clique for which no superset is a clique.

## **More Properties**

- For a tree T with n vertices, the augmented tree G<sub>T</sub> has at most n maximal cliques
- Each maximal clique of G<sub>T</sub> looks like this:
  {v} U N(v)
- There are no other maximal cliques in  $G_T$ .
- If tree T is just a star (one vertex connected to all others), then G<sub>T</sub> is a simple clique
- If G<sub>T</sub> is not a clique, then it has more than one maximal clique, and then T is not a star.

### One more important property

- If (x,y) is an edge of T
  - Then the vertices x and y appear together in exactly two maximal cliques, except if one of them is a leaf
- If one of them is a leaf, then they appear together in exactly one maximal clique

# Properties of Cliques of $G_{\mathsf{T}}$

- Vertex v is present in <= deg(v)+1 maximal cliques</li>
  - Deg(v) is degree of vertex v
- If v has k>0 leaves as neighbors in T, then v is present in exactly deg(v) – k + 1 maximal cliques
- If v has m non-leaves as neighbors in T, then v is in
  - Exactly m + 1 maximal cliques, if v is not a leaf
- If v has no leaves as neighbors in T, then v is in
  - exactly deg(v) + 1 maximal cliques, if v is not a leaf
- If v is a leaf, it is in exactly 1 maximal clique

# Algorithmic Ideas

- **1**. Identify all maximal cliques of  $G_T$
- 2. For each vertex v, compute
  - C[v] = # of maximal cliques of G<sub>T</sub> containing v
- 3. Identify leaves of T: all vertices with C[v] = 1
- 4. Figure out how many non-leaf neighbors each vertex has.
- Figure out pairs of non-leaf vertices connected by an edge (present in exactly 2 max cliques)

## More Properties of leaves of T

- If two leaves x and y are connected to the same non-leaf node, then they appear together in exactly one maximal clique and in no other clique
- If two leaves x and y are not connected to the same non-leaf node, then they never appear together in a maximal clique

# Algorithmic Ideas

- 1. Figure out all leaves of T
- 2. Identify all edges of T connecting non-leaves (skeleton T')
- **3**. Figure out groups of leaves connected to same non-leaf
- 4. Figure out which leaf is connected to which non-leaf:
  - a) Construct skeleton T'
  - b) Construct maximal cliques of T' corresponding to non-leaf
  - c) Each maximal clique A' of T' corresponds to only one maximal clique A of  $G_T$  and to one non-leaf node v.
  - d) Connect all leaf nodes in A to non-leaf node v