

EXAM # 1 ANSWERS

QUESTIONS

Question 1. (30 points)

1. a 2. c 3. a 4. a 5. d 6. b 7. b 8. b 9. b 10. c 11. d 12. d
13. a 14. c 15. a

Grading Criteria: 2 points for each correct answer.

Question 2 (20 points)

Basis: $n = 0$. Then $n = \phi$, the empty set. The only relation from $\{0\}$ to 0 is ϕ , the empty set. But this is not a function since its domain is not $\{0\}$. So, the assertion is vacuously true.

Inductive Step :Assume that the statement is true for n and let $f : (n + 2) = \{0, 1, \dots, n, n + 1\} \longrightarrow \{0, 1, \dots, n\} = (n + 1)$ be one to one. We have two cases, depending on whether n is in the range of f or not.

Case 1: n is not in the range of f . Then, the restriction of f to $n + 1 = \{0, 1, \dots, n\}$ is one to one and its range is a subset of $\{0, 1, \dots, n - 1\} = n$. This contradicts the induction hypothesis.

Case 2: n is in the range of f . Let $i \in (n + 2) = \{0, 1, \dots, n, n + 1\}$ be such that $f(i) = n$. We define a function $g : (n + 2) = \{0, 1, \dots, n, n + 1\} \longrightarrow \{0, 1, \dots, n\} = (n + 1)$ by

$$\begin{aligned} g(i) &= f(n + 1) \\ g(n + 1) &= f(i) \\ g(j) &= f(j) \text{ if } j \neq i \text{ and } j \neq n + 1 \end{aligned}$$

Since f is injective, so is g . The restriction of g to $n + 1$ is one-to-one and its range is a subset of $n = \{0, 1, \dots, n - 1\}$. Again, this contradicts the induction hypothesis.

Grading Criteria:

1. Just trying (listing the cases) : 3 points
 2. Basis: 3 points
 3. The inductive step: 14 points
- 0.3in 3.0 The split by cases: 1 point
- 3.1 Case 1: 3 points
 - 3.2 Case 2: 10 points
 - 3.2.1 The construction of g : 7 points
 - 3.2.2 The rest of the Case 2 proof: 3 points

Question 3 (30 points)

1. a
2. d
3. d
4. b
5. b
6. c
7. b
8. b
9. b
10. c
11. d
12. d
13. a
14. b
15. a

Proof:

We need to prove $\mathcal{P}[F]$.

$\mathcal{P}[\mathcal{F}]$ Every nonempty suffix of F with $n[(, S] = n[), S]$ is a subformula of F .

So we have to show that every S satisfies (*).

(*) if S is a no-empty suffix of F with $n[(, S] = n[), S]$ is a subformula of F .

Case 1: F is an atom.

The only suffixes of F are $S = \lambda$ and $S = F$. If $S = \lambda$, the statement is vacuously true. If $S = F$, S is a formula, so again (*) is true.

Case 2: $F = \neg G$.

The suffixes of F are either suffixes of G or F itself. If S is a suffix of G , then S satisfies (*) by IH. If $S = F$, then (*) is satisfied because S is a formula.

Cases 3,4,5,6: $F = (GCH)$ where C is a binary connective.

The suffixes of F fall into 4 subcase below:

Subcase 3.1: $S = \lambda$.

Then (*) is vacuously true.

Subcase 3.2: $S = J)$ where J is a suffix of H .

$$\begin{aligned}
 n[), S] &= n[), J] + 1 && \text{because } S = J) \\
 &> n[), J] && \text{drop the 1} \\
 &\geq n[(, J] && \text{by Lemma 2 from the class} \\
 &= n[(, J)] \\
 &= n[(, S] && \text{because } S = J) \\
 \text{So, } n[), S] &> n[(, S].
 \end{aligned}$$

For these suffixes S , $(*)$ is vacuously true.

Subcase 3.2: $S = ICH$ where I is a suffix of G .

$$n[), S] = n[), I] + n[), H] + 1 \quad \text{because } S = ICH)$$

$$> n[), I] + n[), H] \quad \text{drop the 1}$$

$$\geq n[(, I] + n[(, H] \quad \text{Lemma 2 from the class notes applied twice}$$

$$= n[(, ICH)$$

$$= n[(, S] \quad \text{because } S = ICH)$$

$$\text{So, } n[), S] > n[(, S].$$

For these suffixes S , $(*)$ is vacuously true.

Subcase 3.4: $S = F$. Since S is a formula, S satisfies $(*)$.

Grading Criteria:

1. Listing the cases: 2 points
2. Case 1: 2 points
3. Case 2: 4 points
4. Cases 3,4,5,6: 12 points
 - 4.1 the split into subcases and Subcase 3.1: 2 points
 - 4.2 Subcase 3.2: 4 points
 - 4.2 Subcase 3.3: 5 points
 - 4.3 Subcase 3.4: 1 point