



Figure 1: The answer to Q2

COT 3420
Spring 2006

EXAM # 2 ANSWERS

Question 1. (39 points)

1. b 2. b 3. a 4. d 5. b 6. c 7. a 8. a 9. d 10. c
11. a 12. c 13. c

Grading Criteria: 3 points for each correct answer.

Question 2. (20 points)

The resolution tree is shown in Figure 1.

Grading Criteria: 1. 20/7 points for each correct derivation step that leads to \square

2. -4 points for each incorrect derivation step
3. -1 point for each missing or misdirected arrow
4. -2 points for finding \square in 9 or more steps

Question 3. (25 points)

Prove the generalized distributivity law,

$$F \vee \bigwedge_{i=1}^n G_i \equiv \bigwedge_{i=1}^n (F \vee G_i).$$

Write your proof on a blank sheet of paper.

Hint: Use induction on n .

Proof: We recall that:

$$\bigwedge_{i=m}^n G_i = \mathbf{T} \text{ if } m > n \text{ (rule 1)}$$

$$\bigwedge_{i=m}^n G_i = G_n \text{ if } m = n \text{ (rule 2)}$$

$$\bigwedge_{i=m}^n G_i = (\bigwedge_{i=m}^{n-1} G_i \wedge G_n) \text{ if } m < n \text{ (rule 3)}$$

Now let's construct the proof.

Case $n = 0$.

$$\text{LHS} = F \vee \bigwedge_{i=1}^0 G_i = F \vee \bigwedge_{i=1}^0 G_i$$

$$= F \vee \mathbf{T} \quad \text{rule 1}$$

$$\equiv \mathbf{T} \quad \text{tautology law}$$

$$\text{RHS} = \bigwedge_{i=1}^0 (F \vee G_i) = \bigwedge_{i=1}^0 (F \vee G_i)$$

$$= \mathbf{T} \quad \text{rule 1}$$

So, both sides are equivalent.

Case $n = 1$.

$$\text{LHS} = F \vee \bigwedge_{i=1}^1 G_i = F \vee \bigwedge_{i=1}^1 G_i$$

$$= F \vee G_1 \quad \text{rule 2}$$

$$\text{RHS} = \bigwedge_{i=1}^1 (F \vee G_i) = \bigwedge_{i=1}^1 (F \vee G_i)$$

$$= F \vee G_1 \quad \text{rule 2}$$

So, both sides are equal.

Case $n = 2$.

$$\text{LHS} = F \vee \bigwedge_{i=1}^2 G_i = F \vee \bigwedge_{i=1}^2 G_i$$

$$= F \vee (\bigwedge_{i=1}^1 G_i \wedge G_2) \quad \text{rule 3}$$

$$= F \vee (G_1 \wedge G_2) \quad \text{rule 2}$$

$$\text{RHS} = \bigwedge_{i=1}^2 (F \vee G_i) = \bigwedge_{i=1}^2 (F \vee G_i)$$

$$= (\bigwedge_{i=1}^1 (F \vee G_i) \wedge (F \vee G_2)) \quad \text{rule 3}$$

$$= ((F \vee G_1) \wedge (F \vee G_2)) \quad \text{rule 2}$$

We prove that LHS \equiv RHS by truth tables.

| F | G_1 | G_2 | $F \vee G_1$ | $F \vee G_2$ | $(F \vee G_1) \wedge (F \vee G_2)$ | $(G_1 \wedge G_2)$ | $F \vee (G_1 \wedge G_2)$ |
|-----|-------|-------|--------------|--------------|------------------------------------|--------------------|---------------------------|
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 | 1 | 0 | 0 | 0 |
| 0 | 1 | 0 | 1 | 0 | 0 | 0 | 0 |
| 0 | 1 | 1 | 1 | 1 | 1 | 1 | 1 |
| 1 | 0 | 0 | 1 | 1 | 1 | 0 | 1 |
| 1 | 0 | 1 | 1 | 1 | 1 | 0 | 1 |
| 1 | 1 | 0 | 1 | 1 | 1 | 0 | 1 |
| 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 |

The columns under $(F \vee G_1) \wedge (F \vee G_2)$ and $F \vee (G_1 \wedge G_2)$ are identical, so they are equivalent. This is the distributivity equivalence.

Inductive step ($n \geq 2$).

Assume (IH).

(IH) $F \vee \bigwedge_{i=1}^n G_i \equiv \bigwedge_{i=1}^n (F \vee G_i)$

We need to show

$F \vee \bigwedge_{i=1}^{n+1} G_i \equiv \bigwedge_{i=1}^{n+1} (F \vee G_i)$

LHS = $F \vee \bigwedge_{i=1}^{n+1} G_i$

= $F \vee (\bigwedge_{i=1}^n G_i \wedge G_{n+1})$ rule 3, because $n + 1 > 1$

$\equiv (F \vee \bigwedge_{i=1}^n G_i) \wedge (F \vee G_{n+1})$ Case $n = 2$

$\equiv \bigwedge_{i=1}^n (F \vee G_i) \wedge (F \vee G_{n+1})$ by IH

= $\bigwedge_{i=1}^{n+1} (F \vee G_i)$ rule 3

= RHS. **Q.E.D.**

Grading Criteria: Case $n = 0$: 4 points; 1 point for the reasons

Case $n = 1$: 4 points; 1 point for the reasons

Case $n = 2$: 9 points; the truth table is worth 5 points; 1.5 points for the reasons

The inductive case: 8 points; the IH 2 points; the derivation 4.5 points; the explanations 1.5 points

Question 4. (25 points)

Show, by structural induction, that the set $S = \{F \wedge \neg G, \mathbf{T}\}$ is adequate where \mathbf{T} is a tautology.

Let \diamond be defined as $F \diamond G = F \wedge \neg G$. We need to show that $S = \{\diamond, \mathbf{T}\}$ is adequate. First, we prove the lemma below.

Lemma: 1. $\neg F \equiv \mathbf{T} \diamond F$

2. $F \vee G \equiv \mathbf{T} \diamond ((\mathbf{T} \diamond F) \diamond G)$

3. $F \wedge G \equiv F \diamond (\mathbf{T} \diamond G)$

4. $F \longrightarrow G \equiv \mathbf{T} \diamond (F \diamond G)$
5. $F \longleftrightarrow G \equiv (\mathbf{T} \diamond (F \diamond G)) \diamond (G \diamond F)$

Proof of the lemma:

1. $\neg F \equiv \mathbf{T} \wedge \neg F$ tautology law
 $= \mathbf{T} \diamond F$ definition of \diamond
2. $F \vee G \equiv \neg\neg(F \vee G)$ double \neg introduction
 $\equiv \neg(\neg F \wedge \neg G)$ DeMorgan's law
 $\equiv \neg(\neg F \diamond G)$ definition of \diamond
 $\equiv \neg((\mathbf{T} \diamond F) \diamond G)$ part 1
 $\equiv \mathbf{T} \diamond ((\mathbf{T} \diamond F) \diamond G)$ part 1
3. $F \wedge G \equiv F \wedge \neg\neg G$ double \neg introduction
 $\equiv F \diamond \neg G$ definition of \diamond
 $\equiv F \diamond (\mathbf{T} \diamond G)$ part 1
4. $F \longrightarrow G \equiv \neg F \vee G$ \longrightarrow -elim
 $\equiv \neg\neg(\neg F \vee G)$ double \neg introduction
 $\equiv \neg(\neg\neg F \wedge \neg G)$ DeMorgan's law
 $\equiv \neg(F \wedge \neg G)$ double \neg elim
 $\equiv \neg(F \diamond G)$ definition of \diamond
 $\equiv \mathbf{T} \diamond (F \diamond G)$ part 1
5. $F \longleftrightarrow G \equiv (F \longrightarrow G) \wedge (G \longrightarrow F)$ \longleftrightarrow -elim
 $\equiv (\neg F \vee G) \wedge (\neg G \vee F)$ \longrightarrow -elim
 $\equiv (\neg F \vee \neg\neg G) \wedge (\neg G \vee \neg\neg F)$ double \neg introduction twice
 $\equiv \neg(F \wedge \neg G) \wedge \neg(G \wedge \neg F)$ DeMorgan's law twice
 $\equiv \neg(F \diamond G) \wedge \neg(G \diamond F)$ definition of \diamond
 $\equiv \neg(F \diamond G) \diamond (G \diamond F)$ definition of \diamond
 $\equiv (\mathbf{T} \diamond (F \diamond G)) \diamond (G \diamond F)$ part 1

Q.E.D. lemma

Now, let us prove that $S = \{\diamond, \mathbf{T}\}$ is adequate by structural induction on F .

Proof:

Case 1: F is an atom. Then F is an S -formula.

Case 2: $F = \neg G$. By IH there is an S -formula G_1 such that $G \equiv G_1$. Then

$$\begin{aligned}
 F &= \neg G \\
 &\equiv \neg G_1 && \text{by IH} \\
 &\equiv \mathbf{T} \diamond G_1 && \text{by Lemma, part 1}
 \end{aligned}$$

The last formula is an S -formula.

Case 3: $F = G \vee H$. By IH there are S -formulas G_1 and H_1 such that $G \equiv G_1$ and $H \equiv H_1$. Then

$$\begin{aligned} F &= G \vee H \\ &\equiv G_1 \vee H_1 && \text{by IH} \\ &\equiv \mathbf{T} \diamond ((\mathbf{T} \diamond G_1) \diamond H_1) && \text{Lemma, part 2} \end{aligned}$$

The last formula is S .

The rest of the cases ($F = G \wedge H$, $F = G \longrightarrow H$, $F = G \longleftrightarrow H$) are solved in a similar fashion, using the corresponding parts of the lemma. **Q.E.D.**

Grading Criteria: 1. Listing the 6 cases: 3 points

2. Case 1: 1 point

3. Case 2: 5 points (the corresponding part of the lemma : 3.5 points, the IH 1, the derivation 0.5 points)

4. Case 3: 4 points (the corresponding part of the lemma : 3 points, the IH 0.5, the derivation 0.5 points)

5. Case 4: 4 points

6. Case 5: 4 points

7. Case 6: 4 points

Question 5. (20 points)

Find a CNF for the formula

$$F = [\neg(A \wedge (B \vee C)) \longleftrightarrow (\neg B \vee \neg(C \vee D))].$$

Solution

$$\begin{aligned} F &= [\neg(A \wedge (B \vee C)) \longleftrightarrow (\neg B \vee \neg(C \vee D))] && \text{line 1} \\ &\equiv [\neg(A \wedge (B \vee C)) \longrightarrow (\neg B \vee \neg(C \vee D))] \wedge [(\neg B \vee \neg(C \vee D)) \longrightarrow \neg(A \wedge (B \vee C))] && \longleftrightarrow\text{-elim, line 2} \\ &\equiv [\neg\neg(A \wedge (B \vee C)) \vee (\neg B \vee \neg(C \vee D))] \wedge [\neg(\neg B \vee \neg(C \vee D)) \vee \neg(A \wedge (B \vee C))] && \longrightarrow\text{-elim, line 3} \\ &\equiv [(A \wedge (B \vee C)) \vee (\neg B \vee \neg(C \wedge \neg D))] \wedge [(\neg\neg B \wedge \neg\neg(C \vee D)) \vee (\neg A \vee \neg(B \vee C))] && \text{double } \neg\text{-elim, De Morgan's 3 times, line 4} \\ &\equiv [(A \wedge (B \vee C)) \vee (\neg B \vee (\neg C \wedge \neg D))] \wedge [(B \wedge (C \vee D)) \vee (\neg A \vee (\neg B \wedge \neg C))] && \text{double } \neg\text{-elim twice, De Morgan's, line 5} \\ &\equiv [(A \wedge (B \vee C)) \vee ((\neg B \vee \neg C) \wedge (\neg B \vee \neg D))] \wedge [(B \wedge (C \vee D)) \vee ((\neg A \vee \neg B) \wedge (\neg A \vee \neg C))] && \text{distributivity twice, line 6} \\ &\equiv (A \vee \neg B \vee \neg C) \wedge (A \vee \neg B \vee \neg D) \wedge (B \vee C \vee \neg B \vee \neg C) \wedge (B \vee C \vee \neg B \vee \neg D) \wedge (B \vee \neg A \vee \neg B) \wedge (B \vee \neg A \vee \neg C) \wedge (C \vee D \vee \neg A \vee \neg B) \wedge (C \vee D \vee \neg A \vee \neg C) && \text{generalized distributivity, line 7} \\ &\equiv (A \vee \neg B \vee \neg C) \wedge (A \vee \neg B \vee \neg D) \wedge (\neg A \vee B \vee \neg C) \wedge (\neg A \vee \neg B \vee C \vee D) && \text{tautology elim, ordering the literals, line 8} \end{aligned}$$

Grading Criteria:

You get credit up to the first line where you made an error.

If the error occurred on line 2 you get 1 point.

If the error occurred on line 3 you get 2 points.

If the error occurred on line 4 you get 4 points.

If the error occurred on line 5 you get 7 points.

If the error occurred on line 6 you get 10 points.

If the error occurred on line 7 you get 13 points.

If the error occurred on line 8 you get 16 points.

If there is no error 19 points.

1 point is for reasons.