

COT 3420  
Section U1  
FALL 2006

EXAM # 1 ANSWERS

**Question 1.** (28 points)

1. d 2. d 3. c 4. a 5. c 6. b 7. a 8. c 9. c 10. a 11. c  
12. a 13. d 14. c

**Grading criteria:** 2 points for each correct answer.

**Question 2.** (25 points)

Do parts a and b.

a. (12 points) Prove that the set of connectives  $S_1 = \{(F \vee \neg G) \wedge H\}$  is not adequate.

**Proof:** The idea is find a property that is true for all  $S_1$  formulas, but false for some formulas in *FORM*. We claim that the truth assignment  $\mathcal{A}$  that assigns 1 to all atoms is a model of all  $S_1$  formulas. We prove it by structural induction on  $S_1$ .

Basis: If  $F$  is an atom  $P_i$ , then  $\mathcal{A}[P_i] = 1$  from the definition of  $\mathcal{A}$ .

Inductive step: Assume that  $F = (G \vee \neg H) \wedge I$ .

We get (1) and (2) by IH on  $G$  and  $H$ .

(1)  $\models_{\mathcal{A}} G$

(2)  $\models_{\mathcal{A}} I$

Since  $G \vee \neg H$  is a consequence of  $G$ , (1) implies (3).

(3)  $\models_{\mathcal{A}} (G \vee \neg H)$

The formula  $(G \vee \neg H) \wedge I$  is a consequence of the set  $\{(G \vee \neg H), I\}$ . So, (3) and (2) yield (4).

(4)  $\models_{\mathcal{A}} ((G \vee \neg H) \wedge I)$ .

**Grading Criteria:**

1. 3 points for guessing an unsatisfiable formula.
2. 2 points for the basis.
3. 7 points for the inductive step.

b. (13 points) Show that the set of connectives  $S_2 = \{(F \vee \neg G) \wedge \neg H\}$  is adequate. You may use the fact that the sets  $\{\neg F, F \wedge G\}$  and  $\{\neg F, F \vee G\}$  are adequate.

We will show that the subset of  $S_2$ ,  $\{(F \vee \neg G) \wedge \neg F\}$  is adequate.

**Proof:**

Let  $\phi[F, G] = (F \vee \neg G) \wedge \neg F$ .

Then,

$\phi[F, F] = (F \vee \neg F) \wedge \neg F$  def of  $\phi$

$\equiv \mathbf{T} \wedge \neg F$  tautology law

$\equiv \neg F$  tautology law

So,

(1)  $\sigma_{\neg} = \phi[F, F]$

Now let compute  $\phi[\phi[G, F], \phi[G, F]]$ .

$\phi[\phi[G, F], \phi[G, F]] \equiv \neg\phi[G, F]$  by (1)

$\equiv \neg((G \vee \neg F) \wedge \neg G)$  def of  $\phi$

$\equiv \neg(G \vee \neg F) \vee \neg\neg G$  De Morgan's law

$\equiv (\neg G \wedge \neg\neg F) \vee G$  De Morgans', double neg elim

$\equiv (\neg G \wedge F) \vee G$  double neg elim

$\equiv (\neg G \vee G) \wedge (F \vee G)$  distribute  $\wedge$  over  $\vee$

$\equiv \mathbf{T} \wedge (F \vee G)$  tautology law

$\equiv F \vee G$  tautology law

So,

(2)  $\sigma_{\vee} = \phi[\phi[G, F], \phi[G, F]]$ .

**Grading Criteria:**

1. 7 points for  $\sigma_{\neg} = \phi[F, F]$

2. 6 points for  $\sigma_{\vee} = \phi[\phi[G, F], \phi[G, F]]$ .

3. If you did not find  $\sigma_{\neg}$  or  $\sigma_{\vee}$  expressed as  $\phi$ -formulas, you cannot get more than 4 points.

**Question 3.** (15 points)

Prove or disprove: If  $F \longrightarrow G$  is a tautology, and  $F$  and  $G$  have no atoms in common, then at least one of  $\neg F$  and  $G$  is a tautology.

First you must write Proof or Disproof and then provide the proof or the counter-example.

**Proof:** We assume that  $G$  is not a tautology and we will show that  $\neg F$  is a tautology.

Since  $G$  is not a tautology, it has a countermodel  $\mathcal{A}$ .

(1)  $\mathcal{A}[G] = 0$

Now let  $\mathcal{B}$  be a truth assignment. We define the truth assignment  $\mathcal{C}$  as follows:

$\mathcal{C}[P_i] = \mathcal{B}[P_i]$  if  $P_i$  is in  $F$

$$\mathcal{C}[P_i] = \mathcal{A}[P_i] \quad \text{if } P_i \text{ is not in } F$$

$\mathcal{C}$  agrees with  $\mathcal{B}$  on  $F$ , so (2) holds.

$$(2) \mathcal{C}[F] = \mathcal{B}[F]$$

Since  $F$  and  $G$  have no atoms in common,  $\mathcal{C}$  agrees with  $\mathcal{A}$  on  $G$ .

$$(3) \mathcal{C}[G] = \mathcal{A}[G]$$

Since  $F \rightarrow G$  is a tautology,

$$(4) \mathcal{C}[F \rightarrow G] = 1.$$

Now,

$$\mathcal{C}[F \rightarrow G] = \mathcal{C}[\neg F \vee G] \quad \rightarrow \text{elimination}$$

$$= \mathcal{C}[\neg F] \boxed{\vee} \mathcal{C}[G] \quad \text{interpretation of } \vee$$

$$= \mathcal{B}[\neg F] \boxed{\vee} \mathcal{C}[G] \quad \text{by (2)}$$

$$= \mathcal{B}[\neg F] \boxed{\vee} \mathcal{A}[G] \quad \text{by (3)}$$

$$= \mathcal{B}[\neg F] \boxed{\vee} 0 \quad \text{by (1)}$$

$$= \mathcal{B}[\neg F] \quad \text{from the tables of } \vee$$

We collect the ends of this sequence of equalities and get (5).

$$(5) \mathcal{C}[F \rightarrow G] = \mathcal{B}[\neg F]$$

From (4) and (5) we obtain (6).

$$(6) \mathcal{B}[\neg F] = 1$$

Since  $\mathcal{B}$  is arbitrary,  $\models \neg F$ .

#### Grading Criteria:

1. If you wrote **Disproof** you get 2 points.
2. If you wrote **Proof** you get 5 points plus the points for the proof ( up to 10).
3. The definition of  $\mathcal{C}$  is worth 6 points, the rest of the proof 4 points.

#### Question 4. (14 points)

Construct a derivation tree of  $\square$  from  $S = \{\{A, B, C\}, \{A, B, \neg C\}, \{A, \neg B, D\}, \{A, \neg B, \neg D\}, \{\neg A, D\}, \{\neg A, \neg D, E\}, \{\neg A, \neg D, \neg E\}\}$ .

**Answer:** The tree is shown in Figure 1.

**Grading Criteria:** 1. 2 points for each correct resolution (up to 7) that leads to  $\square$ .

2. -3 points for each incorret resolution step.
3. -4 points for presenting a derivation sequence instead of a derivation tree.

#### Question 5. (19 points)

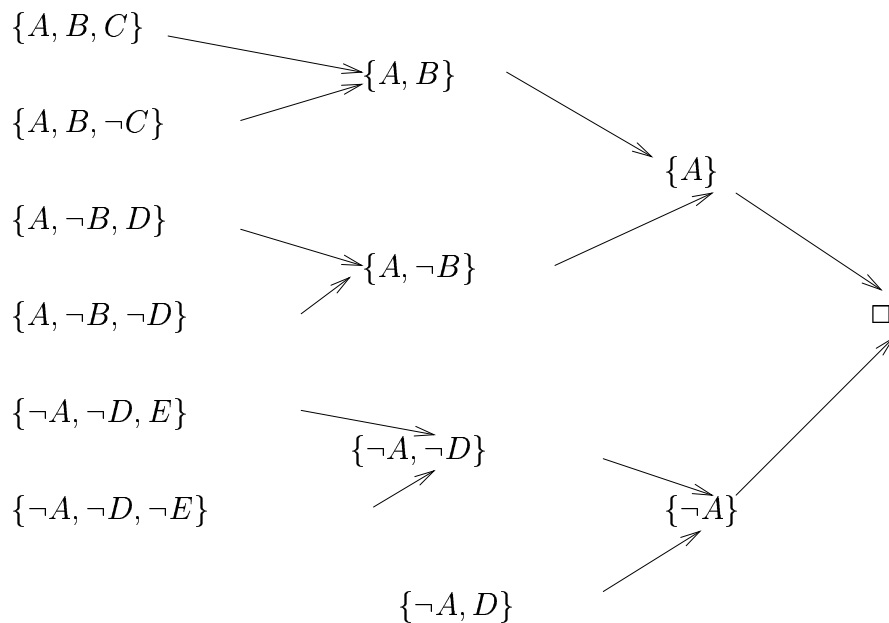


Figure 1: The answer to Question 4

Apply the algorithm given in the book to find a CNF for  $F = \{\neg[A \vee \neg(B \wedge C)] \longrightarrow [B \wedge \neg(C \wedge \neg D)]\}$ .

**Solution**

$$\begin{aligned}
 F &= F_1 = \{\neg[A \vee \neg(B \wedge C)] \longrightarrow [B \wedge \neg(C \wedge \neg D)]\} \\
 &\equiv \neg\neg[A \vee \neg(B \wedge C)] \vee [B \wedge \neg(C \wedge \neg D)] = F_2 \quad \longrightarrow\text{-elim} \\
 &\equiv [A \vee \neg(B \wedge C)] \vee [B \wedge \neg(C \wedge \neg D)] \quad \text{double neg elim} \\
 &\equiv [A \vee \neg B \vee \neg C] \vee [B \wedge (\neg C \vee \neg\neg D)] \quad \text{DeMorgan's law, twice} \\
 &\equiv [A \vee \neg B \vee \neg C] \vee [B \wedge (\neg C \vee D)] = F_3 \quad \text{double neg elim} \\
 &\equiv (A \vee \neg B \vee \neg C \vee B) \wedge (A \vee \neg B \vee \neg C \vee \neg C \vee D) = F_4 \quad \text{distributivity} \\
 &\equiv (A \vee \neg B \vee \neg C \vee D) = F_5 \quad \text{tautology law, idempotency}
 \end{aligned}$$

**Grading Criteria:** 1. You get credit up to the first line that contains an error.

- 1 correct line: 2 point
- 2 correct lines: 4 points
- 3 correct lines: 6 points
- 4 correct lines: 9 points
- 5 correct lines: 12 points
- 6 correct lines: 15 points

7 correct lines: 19 points

2. If you did not list the reasons, you loose 3 points.